

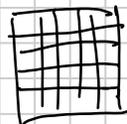
Large-scale matrix equations

Note Title

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Usually, large-scale control systems have:

- large, sparse $A \in \mathbb{C}^{n \times n}$ (from discretization)
- $B \in \mathbb{C}^{n \times m}$ tall thin, $m \ll n$
few controllable degrees of freedom



Simplest possible large-scale equation: Lyapunov equation

$$AX + XA^* + \boxed{bb^*} = 0 \quad b \in \mathbb{C}^n$$

$\exists Q$ has rank-1

If you know how to solve equations with rank-1 RHS,
you can also solve

$$AX + XA^* + b_1 b_1^* + b_2 b_2^* = 0$$

has solution $X_1 + X_2$, where

$$\begin{aligned} AX_1 + X_1 A^* + b_1 b_1^* &= 0, \\ AX_2 + X_2 A^* + b_2 b_2^* &= 0. \end{aligned}$$

Also, if we know how to solve Lyapunov equations, we can
solve large-scale Riccati equations with the Newton method.

We shall also assume $\lambda(A) \subset \text{LHP}$.

Problem: solution X is full, in general

Solution: In most cases, X is well-approximated by a

low-rank matrix, because it has rapidly decaying sing. values.

$$X \approx ZZ^* \quad Z \in \mathbb{C}^{n \times k} \quad k \ll n.$$

Our algorithms will compute Z , not X .

Idea: transform the Lyapunov equation into a (discrete-time)

Stein equation ^(LE) $AX + XA^* + bb^* = 0 \Rightarrow X - \hat{A}X\hat{A}^* = \hat{b}\hat{b}^*$ ^(SE)

Theorem: let $\tau > 0$, A with $\Lambda(A) \subset \text{LHP}$, Then,

X solves the LE if and only if it solves

the Stein equation SE with

$$\hat{A} = (A - \tau I)^{-1}(A + \tau I), \quad \hat{b} = \sqrt{2\tau} (A - \tau I)^{-1} b$$

and $\rho(\hat{A}) < 1$.

Proof: note that $\Lambda(A - \tau I) \subset \text{LHP}$.

$$X - (A - \tau I)^{-1}(A + \tau I)X(A + \tau I)^*(A - \tau I)^{-*} = 2\tau (A - \tau I)^{-1}bb^*(A - \tau I)^{-*}$$

$$(A - \tau I)X(A - \tau I)^* - (A + \tau I)X(A + \tau I)^* = 2\tau bb^* \quad \Leftrightarrow$$

~~$$AXA^* - \tau XA^* - AX\tau + \tau^2 X - AXA^* - \tau XA^* - AX\tau - \tau^2 X = 2\tau bb^*$$~~

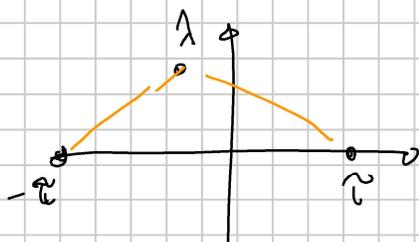
$$\Leftrightarrow 2\tau (AX + XA^* + bb^*) = 0.$$

If $\lambda \in \Lambda(A)$, the corresponding eigenvalue of \hat{A} is $\frac{\lambda + \tau}{\lambda - \tau}$

$$\Lambda(\hat{A}) = \left\{ \frac{\lambda + \tau}{\lambda - \tau}, \text{ where } \lambda \in \Lambda(A) \right\}$$

Since $\lambda \in \text{LHP}$,

$$\text{dist}(\lambda, -\tau) < \text{dist}(\lambda, \tau)$$



i.e. $\frac{|\lambda + \tau|}{|\lambda - \tau|} < 1$, hence $\rho(\hat{A}) < 1$. \square

$$X - \hat{A}X\hat{A}^* = \hat{b}\hat{b}^*$$

$$\rho(\hat{A}) < 1$$

We solve it via

$$\begin{cases} X_0 = 0 \\ X_{k+1} = \hat{b}\hat{b}^x + \hat{A}X_k\hat{A}^* \end{cases} \rightarrow \text{Smith's method}$$

$$X_1 = \hat{b}\hat{b}^x \quad X_2 = \hat{b}\hat{b}^x + \hat{A}\hat{b}\hat{b}^x\hat{A}^*$$

$$X_3 = \hat{b}\hat{b}^x + \hat{A}\hat{b}\hat{b}^x\hat{A}^* + \hat{A}^2\hat{b}\hat{b}^x(\hat{A}^*)^2$$

$$X_k = \sum_{i=0}^{k-1} \hat{A}^i \hat{b}\hat{b}^x (\hat{A}^*)^i$$

In terms of low-rank factors:

$$X_k = Z_k Z_k^*, \text{ with } Z_k = \begin{bmatrix} \hat{b} & \hat{A}\hat{b} & \hat{A}^2\hat{b} & \dots & \hat{A}^{(k-1)}\hat{b} \end{bmatrix} \in \mathbb{C}^{n \times k}$$

$$\begin{cases} v_0 = \hat{b} & k=0,1,2,\dots \\ v_{k+1} = \hat{A}v_k & Z_k = [v_0 \ v_1 \ \dots \ v_{k-1}] \end{cases}$$

$$v_{k+1} = \underbrace{(A - \tau I)^{-1}} (A + \tau I) v_k \quad v_0 = \sqrt{2\tau} \underbrace{(A - \tau I)^{-1}} \hat{b}$$

We only need to solve sparse linear systems!

$$v_{k+1} = (A - \tau I)^{-1} (A - \tau I + 2\tau I) v_k = v_k + 2\tau (A - \tau I)^{-1} v_k$$

$$\begin{aligned} Z_k Z_k^* &= \begin{bmatrix} \hat{b} & \hat{A}\hat{b} & \hat{A}^2\hat{b} & \dots & \hat{A}^{k-1}\hat{b} \end{bmatrix} \begin{bmatrix} \hat{b}^x \\ \hat{b}^x \hat{A}^x \\ \vdots \\ \hat{b}^x (\hat{A}^x)^{k-1} \end{bmatrix} \\ &= \hat{b}\hat{b}^x + \hat{A}\hat{b}\hat{b}^x\hat{A}^x + \dots + \hat{A}^{k-1}\hat{b}\hat{b}^x(\hat{A}^x)^{k-1} \end{aligned}$$

Note that $\|\hat{A}^k\| \sim \rho(\hat{A})^k$, so the sum converges.

X_k converges to a fixed point $X_* = \hat{b}\hat{b}^* + \hat{A}X_*\hat{A}^*$,
 which must be the exact solution of the equation

$$X_* = \sum_{i=0}^{\infty} (\hat{A}^i) \hat{b}\hat{b}^* (\hat{A}^*)^i$$

$$\begin{aligned} X_* - \hat{z}_k \hat{z}_k^* &= \sum_{i=k}^{\infty} \hat{A}^i \hat{b}\hat{b}^* (\hat{A}^*)^i = \hat{A}^k \left(\sum_{i=0}^{\infty} \hat{A}^i \hat{b}\hat{b}^* (\hat{A}^*)^i \right) (\hat{A}^*)^k \\ &= \hat{A}^k X_* (\hat{A}^*)^k \end{aligned}$$

$$\|X_* - \hat{z}_k \hat{z}_k^*\| \leq \|\hat{A}^k\|^2 \|X_*\| \sim \rho(\hat{A})^{2k} \quad \text{converges to } 0, \text{ since } \rho(\hat{A}) < 1$$

Convergence rate:

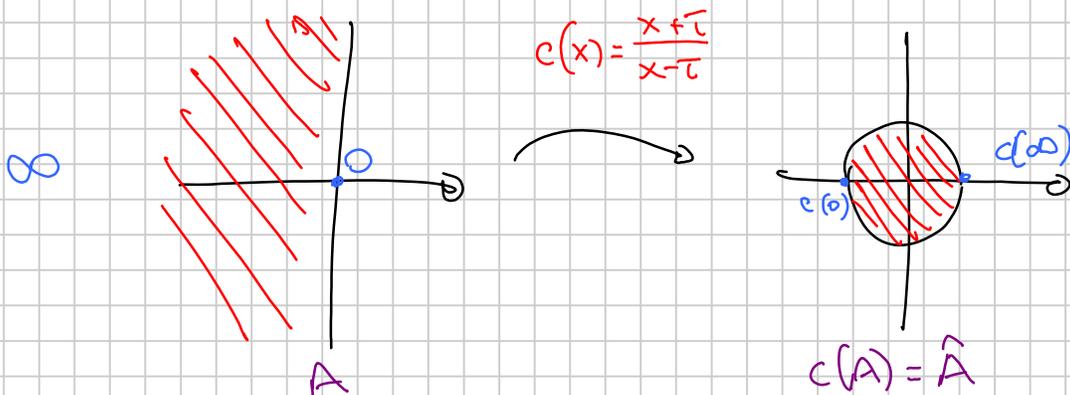
$$\max_{\lambda \in \Lambda(A)} \left| \frac{\lambda + \tau}{\lambda - \tau} \right|^2 < 1$$

If $\tau \gg |\lambda|$, $\frac{\lambda + \tau}{\lambda - \tau} \approx -1 \rightarrow$ slow convergence

If $\tau \ll |\lambda|$, $\frac{\lambda + \tau}{\lambda - \tau} \approx 1 \rightarrow$ slow convergence.

If A has all eigenvalues close, $\tau \approx \lambda$ is best

If A has both small and large eigenvalues, no luck.



Idea: change τ at every iteration

$$A \quad \hat{A}_k = (A - \tau_k I)^{-1} (A + \tau_k I) \quad \hat{b}_k = \sqrt{2\tau_k} (A - \tau_k I)^{-1} b$$

$$X \text{ solves } X = \hat{A}_k X \hat{A}_k^* + \hat{b}_k \hat{b}_k^* \text{ for each } k.$$

$\tau_1, \tau_2, \tau_3, \dots$

$$X_k = \hat{A}_k X_{k-1} \hat{A}_k^* + \hat{b}_k \hat{b}_k^* \quad X_0 = 0$$

$$X_1 = \hat{b}_1 \hat{b}_1^*$$

$$X_2 = \hat{b}_2 \hat{b}_2^* + \hat{A}_2 \hat{b}_1 \hat{b}_1^* \hat{A}_2^*$$

$$X_3 = \hat{b}_3 \hat{b}_3^* + \hat{A}_3 \hat{b}_2 \hat{b}_2^* \hat{A}_3^* + \hat{A}_3 \hat{A}_2 \hat{b}_1 \hat{b}_1^* \hat{A}_2^* \hat{A}_3^*$$

\vdots

\vdots

\vdots

$$Z_k = \left[\hat{b}_k \quad \hat{A}_k \hat{b}_{k-1} \quad \dots \quad \hat{A}_k \hat{A}_{k-1} \dots \hat{A}_2 \hat{b}_1 \right]$$

$$\boxed{\sqrt{2\tau_k} (A - \tau_k I)^{-1} b}$$

v_0

$$\boxed{(A - \tau_k I)^{-1} (A + \tau_k I) (A - \tau_{k-1} I)^{-1} b \sqrt{2\tau_{k-1}}}$$

$$v_1 = (A - \tau_{k-1} I)^{-1} (A + \tau_k I) (A - \tau_k I)^{-1} b \sqrt{2\tau_{k-1}}$$

$$= (A - \tau_{k-1} I)^{-1} (A + \tau_k I) v_0 \frac{\sqrt{2\tau_{k+1}}}{\sqrt{2\tau_k}}$$

ADI with multiple shifts:

$$\alpha_0 = \tau_k \quad \alpha_1 = \tau_{k-1} \quad \alpha_2 = \tau_{k-2}, \dots$$

$$v_0 = \sqrt{2\alpha_0} (A - \alpha_0 I)^{-1} b$$

$$v_j = \frac{\sqrt{2\alpha_j}}{\sqrt{2\alpha_{j-1}}} (v_{j-1} + (\alpha_{j-1} + \alpha_j) (A - \alpha_j I)^{-1} v_{j-1})$$

Computing the residual:

$$\begin{cases} X_k = \hat{A}_k X_{k-1} \hat{A}_k^* + \hat{b}_k \hat{b}_k^* \\ X_* = \hat{A}_k X_* \hat{A}_k^* + \hat{b}_k \hat{b}_k^* \end{cases}$$

$$X_k - X_* = \hat{A}_k (X_{k-1} - X_*) \hat{A}_k^*$$

$$X_k - X_* = \underbrace{\hat{A}_k \hat{A}_{k-1} \dots \hat{A}_1}_{r_k(A)} (X_0 - X_*) \underbrace{\hat{A}_1^* \hat{A}_2^* \dots \hat{A}_k^*}_{r_k(A)^*}$$

$$r_k(\lambda) = \frac{\lambda + \tau_k}{\lambda - \tau_k} \cdot \frac{\lambda + \tau_{k-1}}{\lambda - \tau_{k-1}} \cdot \dots \cdot \frac{\lambda + \tau_1}{\lambda - \tau_1}$$

$$\|X_k - X_*\| \leq \|r_k(A)\|^2 \cdot \|X_0 - X_*\|$$

We want $\|r_k(A)\|$ as small as possible for convergence:

‡ If A symmetric,

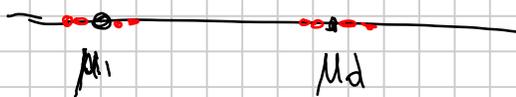
$$\|r_k(A)\| = \max_{\lambda \in \Lambda(A)} \text{eigenvalue} = \max_{\lambda \in \Lambda(A)} \left| \frac{\lambda + \tau_k}{\lambda - \tau_k} \cdot \frac{\lambda + \tau_{k-1}}{\lambda - \tau_{k-1}} \cdot \dots \cdot \frac{\lambda + \tau_1}{\lambda - \tau_1} \right|$$

‡ If A has only a few distinct eigenvalues,

$$\lambda_1, \dots, \lambda_d,$$

We can choose $\tau_1 = -\lambda_1, \dots, \tau_d = -\lambda_d$, and after d steps we have exact convergence, $r_d(A) = 0$

‡ If A has d clusters of eigenvalues centered in μ_1, \dots, μ_d ,



we choose $\tau_i = -\mu_i$

We can formulate an optimization problem for each new shift:

Wachspress shifts

$$\alpha_k = \arg \min_{\lambda \in \Lambda(A)} \max_{\lambda \in \Lambda(A)} \left| \frac{\lambda + \alpha_k}{\lambda - \alpha_k} \right| \cdot \left| \frac{\lambda + \alpha_{k-1}}{\lambda - \alpha_{k-1}} \right| \cdots \left| \frac{\lambda + \alpha_1}{\lambda - \alpha_1} \right|$$

To solve this, we would require the whole $\Lambda(A)$, and this is unfeasible

Workaround: compute approximations of the larger and smaller eigenvalues of A with Arnoldi, and then use them instead of $\Lambda(A)$ to solve the problem.

More theoretical solution: assume $\Lambda(A) \subseteq [a, b]$.

Then, we look for polynomials of degree k that achieve

$$\alpha(k) = \min_{P \in \mathcal{P}_k} \max_{\lambda \in [a, b]} \left| \frac{P(\lambda)}{P(-\lambda)} \right| \quad (\text{Zolotar problem})$$

People know solutions that achieve the optimal decrease,

$$\alpha(k) \sim r^k \quad \text{for a given } r < 1, \text{ "logarithmic capacity" of } [a, b].$$

$$Z_k = \begin{bmatrix} \hat{b}_k & \hat{A}_k \hat{b}_{k-1} & \cdots & \hat{A}_k \hat{A}_{k-1} \cdots \hat{A}_2 \hat{b}_1 \end{bmatrix}$$

$$\hat{A}_k = (A - \tau_k I)^{-1} (A + \tau_k I)$$

$$\hat{b}_k = \sqrt{2\tau_k} (A - \tau_k I)^{-1} b$$

All columns of Z_k belong to

$$\text{span} \left((A - \tau_1 I)^{-1} b, (A - \tau_2 I)^{-1} b, \dots, (A - \tau_k I)^{-1} b \right)$$

$$= \text{span} \left(r(A)b \text{ s.t. } r(z) \text{ is a rational function} \right.$$

$$\left. \text{with denominator } (z - \tau_1)(z - \tau_2) \dots (z - \tau_k) \right)$$

= rational Krylov subspace

Idea: first compute a basis for the space V_k ,
then choose an approximation of the solution of
the form

$$X_k = V_k Y_k V_k^*$$

$$\boxed{\quad} \square \square \quad Y_k \in \mathbb{C}^{k \times k}$$

Projected equation:

$$V_k^* A V_k Y_k V_k^* + V_k^* V_k Y_k V_k^* A^* V_k + V_k^* b b^* V_k = 0$$

$$A_k Y_k + Y_k A_k^* + (V_k^* b)(b^* V_k) = 0. \quad \text{A } k \times k \text{ equation}$$

(rational Krylov algorithms)

$$A_k = V_k^* A V_k$$