## Large-scale methods for Lyapunov equations

We give a hint of the methods used for large-scale equations.

We focus on Lyapunov equations,  $AX + XA^* + BB^* = 0$ . (Then we can solve CAREs using Newton's method, for instance.)

Assumptions: A large and sparse with  $\Lambda(A) \subset LHP$ .  $B \in \mathbb{R}^{n \times m}$ , with  $m \ll n$ .

Actually, we may suppose  $B = b \in \mathbb{R}^n$  without loss of generality: a rank-*m* matrix is the sum of *m* rank-1 matrices, and the equation is linear.

Assume A symmetric, normal or 'almost normal'. The algorithms often work for generic A, but the analysis works better for normal matrices.

# ADI (alternating-direction implicit) iteration

Let  $p_k < 0 \in \mathbb{R}$ , so that  $A + p_k I$  is never singular (and  $A - p_k I$  may be singular if  $p_k$  is an eigenvalue of A).

Rewrite the Lyapunov equation as

$$(A + p_k I)X + X(A^* - p_k I) + bb^* = 0.$$

and formulate as a fixed-point equation (starting from  $X_0 = 0$ )

$$X_{k+1} = (A + p_k I)^{-1} (-bb^* - X_k (A^* - p_k I)).$$

Ugly, breaks symmetry: to restore it, make two steps with same  $p_k$ :

$$\begin{aligned} X_{k+\frac{1}{2}} &= (A + p_k I)^{-1} (-bb^* - X_k (A^* - p_k I)), \\ X_{k+1} &= (-bb^* - (A - \overline{p_k} I) X_{k+\frac{1}{2}}) (A^* + \overline{p_k} I)^{-1} \\ &= -2 \operatorname{Re}(p_k) (A + p_k I)^{-1} bb^* (A + p_k I)^{-*} + f_{p_k} (A) X_k f_{p_k} (A)^*, \end{aligned}$$

with  $f_{p_k}(x) = \frac{x - \overline{p_k}}{x + p_k} = 1 - 2\operatorname{Re}(p_k)\frac{1}{x + p_k}$ .

### Low-rank ADI

$$X_{k+1} = -2 \operatorname{Re}(p_k) (A + p_k I)^{-1} b b^* (A + p_k I)^{-*} + f_{p_k}(A) X_k f_{p_k}(A)^*$$

Can be rewritten in terms of the 'low-rank factor' of  $X_k = Z_k Z_k^*$ :

$$Z_{k+1} = \begin{bmatrix} \sqrt{-2\operatorname{Re}(p_k)}(A+p_kI)^{-1}b & f_{p_k}(A)Z_k \end{bmatrix}.$$

or in a more efficient form with only one inversion at each step:

$$Z_k = \begin{bmatrix} v_1 & v_2 & \dots & v_k \end{bmatrix}, \quad v_{k+1} = \alpha_k v_k + \beta_k (A + p_k I)^{-1} v_k.$$

(easier to see by looking at the first steps).

## ADI: convergence

Convergence depends on the choices of  $p_k$ . Intuitively: good if  $A - p_k I$  is small and  $A + p_k I$  is large. Suggests taking  $p_k$  as (some of) the eigenvalues of A.

More formally:

$$egin{aligned} &X_{k+1}-X_*=f_{p_k}(A)(X_k-X_*)f_{p_k}(A)^*=\cdots=g(A)(X_0-X_*)g(A)^*, \ & ext{where }g(x)=\prod_{i=0}^krac{x-\overline{p_i}}{x+p_i}. \ & ext{If }A=V\Lambda V^{-1}, ext{ then} \ &\|g(A)\|\leq\kappa(V)\max_{\lambda\in\Lambda(A)}\prod_{i=0}^krac{|\lambda-\overline{p_i}|}{|\lambda+p_i|}. \end{aligned}$$

How to choose  $p_k$ 's that make this small? Easy if A has few / clustered eigenvalues.

#### ADI convergence

$$\eta_k = \min_{p_0, \dots, p_k} \max_{\lambda \in \Lambda(A)} \prod_{i=0}^{k-1} \frac{|\lambda - \overline{p_i}|}{|\lambda + p_i|}.$$

In general, tricky approximation theory problem. Typical approach: find an enclosing region for the eigenvalues of A (for instance, if  $A = A^*$ , all eigenvalues are in  $[\lambda_{\min}, \lambda_{\max}]$ ); look for a polynomial that is 'small' on this region and 'large' on  $[-\lambda_{\max}, -\lambda_{\min}]$ .

In many cases,  $\eta_k$  decays with k — exponentially, or anyway fast.

Consequence The solution X has decaying singular values / low numerical rank.

(Good thing, because otherwise the problem would be hopeless: X is full.)

# Residual computation

For 
$$X_k = Z_k Z_k^*$$
, with  $Z_k \in \mathbb{R}^{n \times k}$ , we have  
 $AZ_k Z_k^* + Z_k Z_k^* A^* + BB^* = \begin{bmatrix} Z_k & AZ_k & B \end{bmatrix} \begin{bmatrix} 0 & I & 0 \\ I & 0 & 0 \\ 0 & 0 & I \end{bmatrix} \begin{bmatrix} Z_k & AZ_k & B \end{bmatrix}^*$ 

Using QR or SVD of the tall thin  $\begin{bmatrix} Z_k & AZ_k & B \end{bmatrix}$ , we can compute residual norms in  $O(nk^2)$ .

# Rational Arnoldi

The computed  $Z_k$  has columns of the form r(A)b, where r(x) = q(x)/p(x), with denominator  $p(x) = (x - p_0)(x - p_1) \dots (x - p_{k-1})$ .

Definition

The rational Krylov subspace with poles  $p_0, p_1, \ldots, p_{k-1}$  is

$$R_k(A,b) = \{p(A)^{-1}q(A)b \colon \deg q < k\} = p(A)^{-1}K_k(A,b).$$

An orthogonal basis can be computed with a variant of Arnoldi: at each step, multiply the last vector  $v_k$  by  $(A - p_k I)^{-1}$  and orthogonalize against the previous ones.

Plan First compute this subspace, then solve the projected equation.

## Solving projected equations

Given an orthonormal basis  $U_k$  of  $R_k(A, b)$ :

1. Set  $X_k = U_k Y_k U_k^*$ ;

2. Assume 'orthogonal residual':  $U_k^*(AX_k + X_kA^* + BB^*)U_k = 0$ . Produces a projected Lyapunov equation

$$(U_k^* A U_k) Y + Y (U_k^* A U_k)^* + U_k^* B B^* U_k = 0.$$

Difficulty 1 Even if A stable,  $U_k^*AU_k$  is not necessarily so. Difficulty 2 (main one, common to ADI): good pole selection. Pole selection can be critical for convergence. No good general strategies. Usually one tries some extremal eigenvalues of A and  $A^{-1}$  as  $p_k$ .