

Derivatives

1) Computer algebra system

2)  $\frac{f(x+h) - f(x)}{h}$  (error  $\approx 10^{-8}$ ) (numerical derivatives)

3) Write "matrix-friendly" code + evaluate in  $\begin{bmatrix} x & x' & \dots & 1 \\ 0 & x' & \dots & x \end{bmatrix}$

to get  $\begin{bmatrix} f(x) & f'(x) & \dots & \vdots \\ \vdots & \vdots & \vdots & \vdots \\ 0 & \vdots & f(x) & \vdots \end{bmatrix}$  (use dimension  $n$  to get  $f(x), \dots, f^{(n-1)}(x)$ )

variables, e.g.  $a \iff A = \begin{bmatrix} a & a' & a'' & \dots & \frac{a^{(n-1)}}{(n-1)!} \\ & \vdots & \vdots & \vdots & \vdots \\ & & 0 & \vdots & a' \\ & & & & a \end{bmatrix}$  (Triangular Toeplitz)

This could be described even without matrices:

$$c = a * b \quad \left| \quad \begin{array}{cc} [a \ a' \ a''] & [b \ b' \ b''] \\ c = [a \ a' \ a''] * [b \ b' \ b''] = \\ & = [ab \ ab' + a'b \ ab'' + 2a'b' + a''b] \end{array} \right.$$

Special case:  $n=2$  (first derivative only)

"dual numbers": replace  $a$  with  $a + \epsilon a'$

Use algebra +  $\epsilon^2 = 0$

$$c = ab \quad (a + \epsilon a')(b + \epsilon b') \\ c + \epsilon c' = ab + \epsilon(ab' + a'b) + \cancel{\epsilon^2 a'b}$$

(The input variable  $x$  becomes  $x + \epsilon \cdot 1$ )

Ways to see this:

► operations in  $\mathbb{R}[\epsilon] / (\epsilon^2)$

►  $\epsilon = \begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix}$

► " $\epsilon$  is infinitesimal"

"Complex step differentiation":

Suppose you have  $f$  holomorphic / analytic

You are interested in  $f(x) \in \mathbb{R}$   $x \in \mathbb{R}$  and  $f'(x)$

You have code that computes  $f$  and works also with complex input

$$f(x+ih) = \underbrace{f(x)} + ih \underbrace{f'(x)} - \frac{h^2}{2} \underbrace{f''(x)} + O(h^3)$$

$$\frac{\text{Im}(f(x+ih))}{h} = f'(x) + O(h^2) \quad \left( \begin{array}{l} \text{one order of precision} \\ \text{better than } \frac{f(x+h)-f(x)}{h} = \\ = f'(x) + O(h) \end{array} \right)$$

In the numerator: error approx.  $u \cdot f'(x)$

$$\text{In total: error } O\left(\frac{u \cdot f'(x)}{h} + h^2\right)$$

$$h \approx u^{1/3} \text{ gives } O\left(f'(x)u^{2/3} + u^{2/3}\right) \quad u \sim 10^{-16}$$
$$\sim 10^{-2/3 \cdot 16} \sim 10^{-11}$$

$a \gg b$

$$f(x) = (a+ib)^3 = \underbrace{a^3 + 3a^2ib - 3ab^2 - ib^3}$$

$x = a+ib$

error in computing  $\text{Im} f(x)$   
 $3a^2b - b^3$  is proportional  
to  $\text{Im} f(x) \ll f(x)$

Previous (exact) techniques: automatic differentiation.

What we described (extending operations to derivatives)  
is forward mode of automatic differentiation

There is also reverse mode ("back-propagation" in machine learning)

General idea: first you run your code to compute  $f(x) = y$   
then you go through instruction backwards and compute  
for each variable  $a$   $\frac{\partial y}{\partial a}$ .

(We will not see details here)

Remark: In case you have to compute  $Jf \in \mathbb{R}^{m \times n}$  for

a function  $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$

forward mode is faster if  $n \ll m$ ,

reverse mode is faster for  $n \gg m$ . (machine learning)

## Methods for specific matrix functions

First one:  $\expm(A) = I + A + \frac{A^2}{2} + \frac{A^3}{3!} + \dots$

The solution of the vector-valued ODE

$$\frac{d}{dt} v(t) = Av(t) \quad v(0) = v_0$$

$$v: \mathbb{R} \rightarrow \mathbb{R}^n \\ A \in \mathbb{R}^{n \times n}$$

is  $v(t) = \expm(At) \cdot v_0$

(Proof: we can differentiate term by term)

$$v(t) = \left( I + tA + t^2 \frac{A^2}{2} + t^3 \frac{A^3}{3!} + \dots \right) v_0$$

$$\frac{d}{dt} v(t) = \left( 0 + A + 2t \frac{A^2}{2} + 3t^2 \frac{A^3}{3!} + \dots \right) v_0 = A \left( I + tA + t^2 \frac{A^2}{2} + \dots \right) v_0 \\ = Av(t).$$

We have seen a few problematic methods:

$$A = VJV^{-1} \text{ (Jordan form) } \leadsto \kappa(V) \text{ big, not so good}$$

Series  $I + A + \frac{A^2}{2} + \dots \leadsto$  problems, intermediate growth...

Moler, Van Loan '78 o '03 "Nineteen dubious ways to compute the matrix exponential".

Problem with Maclaurin series: there is large intermediate growth in matrix powers.

$$\|A^n\| \sim \rho(A)^n \quad \|A^n\|^{1/n} \rightarrow \rho(A)$$

Even if  $\rho(A) < 1$  and so  $\|A^n\| \rightarrow 0$ , the intermediate powers may grow a lot.

Ex:

$$A = \begin{bmatrix} 0 & 1000 & 0 & 0 \\ 0 & 0 & 100 & 0 \\ 0 & 0 & 0 & 10 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A^2 = \begin{bmatrix} 0 & 0 & 10000 & 0 \\ 0 & 0 & 0 & 1000 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A^3 = \begin{bmatrix} 0 & 0 & 0 & 100000 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} \quad A^4 = 0$$

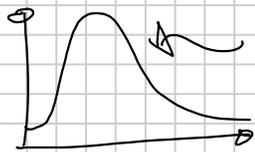
$$\rho(A) = 0$$

Typical behavior for a non-normal matrix

(normal means  $A = QDQ^*$ , with  $Q$  unitary and  $D$  diagonal)

(for a normal matrix,  $A^n = QD^nQ^*$   $\|A^n\| = \|A\|^n$ )

For many matrices, the summands in  $\sum \frac{A^k}{k!} = \expm(A)$  have a "hump"



and so does  $\expm(tA)$

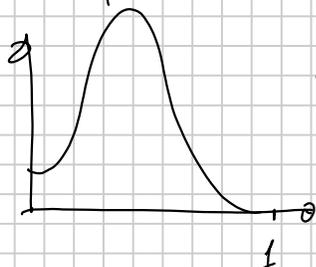
(for matrices that are far from normal)

Other dubious methods: use a method to discretize ODEs to solve

$$\frac{d}{dt}V(t) = AV \quad V(0) = I.$$

$$V(t) = \exp(At)$$

Running an ODE solver up to  $t=1$  should get  $V(1) = \exp(A)$

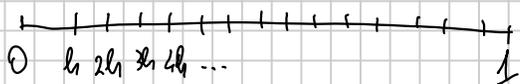


but one might have intermediate growth  $\Rightarrow$  large errors.

Explicit Euler method

$$AV = \frac{d}{dt}V(t) \approx \frac{V(t+h) - V(t)}{h}$$

$V(0) = V_0$  Discretize  $[0,1]$  into  $n$  points with spacing  $h = \frac{1}{n}$



$$\frac{V(t+h) - V(t)}{h} \approx AV(t)$$

$$V(h) = (I + hA)V_0$$

$$V(h) = V(0) + hAV(0) =$$

$$V(2h) = V(h) + hAV(h)$$

⋮

$$V(2h) = (I + hA)V(h) = (I + hA)^2 V_0$$

$$V(nh) = (I + hA)^n V_0 \quad V(1) \approx \left(I + \frac{1}{n}A\right)^n$$

Padé approximants: are a variant of Maclaurin series

Idea: fix degrees  $P, Q$

We would like to find

$$\exp(x) \approx \frac{N(x)}{D(x)} \quad \begin{array}{l} \text{+ degree } p \\ \text{+ degree } q \end{array}$$

in the sense that  $\exp(x) = \frac{N(x)}{D(x)} + O(x^m)$  with  $m$  as large as possible

eg.  $p=q=2$

$$\exp(x) \approx \frac{ax^2+bx+c}{dx^2+ex+1} = 1+x + \frac{x^2}{2} + \frac{x^3}{3!} + \frac{x^4}{4!} + O(x^5)$$

take  $d=1$  as normalization

$p=2, q=1$

$$\frac{ax^2+bx+c}{1+ex} = (ax^2+bx+c)(1-ex+e^2x^2-e^3x^3+\dots)$$

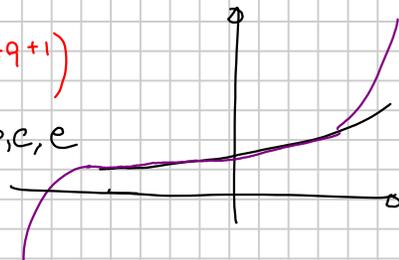
$p+1$  unknowns  
 $q+1$  unknowns  
but you normalize by setting one to 1

$$= c + (b-e)x + (a+ce^2-be)x^2 + (-a+be^2-ce^3)x^3 + O(x^4)$$

$$\stackrel{!}{=} 1 + x + \frac{x^2}{2} + \frac{x^3}{3!} + O(x^4)$$

$$(*) \begin{cases} 1 = c \\ 1 = b - e \\ \frac{1}{2} = a + ce^2 - be \\ \frac{1}{3!} = -a + be^2 - ce^3 \end{cases}$$

solve in  $a, b, c, e$



(Generalization of Maclaurin expansion, which has degree  $\infty$  in the denominator.)

Pade approximations approximate the exponential better than Maclaurin expansion for small  $x$

(There is a trick to convert  $(*)$  into a linear system — can you see it?)