

Arnoldi method

Note Title

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$$K_n(A, b) = \text{span}(b, Ab, A^2b, \dots, A^{n-1}b)$$

(If at step n $A^n b \in \text{span}(b, Ab, \dots, A^{n-1}b)$ for the first time, then $\alpha_{n+1, n} = 0$ and $AV_n = V_n A_n$ is an invariant subspace relation and $f(A)V_n = V_n f(A_n) \forall f$)

$$\begin{matrix} m \times m & m \times n & m \times n & n \times n \\ \boxed{A} & \boxed{V_n} & = & \boxed{A_n} \end{matrix}$$

Arnoldi produces V_n orthonormal basis of $K_n(A, b)$ and $A_n = V_n^* A V_n$

Theorem:

$$c = \underbrace{V_n}_{\substack{\text{coordinates} \\ \text{of } b \in K_n \\ \text{in basis } V_n}} \underbrace{f(A_n)}_{\substack{\text{A restricted \& projected} \\ \text{onto } K_n}} \underbrace{V_n^* b}_{\substack{\text{coordinates} \\ \text{of } b \in K_n \\ \text{in basis } V_n}} = f(A)b$$

when f is a polynomial of degree $< n$

Coordinates in basis V_n

Even if f is not a polynomial, c approximates well $f(A)b$

Heuristic reason: $f(A_n) \approx \tilde{p}(A_n)$, where \tilde{p} is the interpolation polynomial of f in $\Lambda(A_n)$.

\tilde{p} has degree at most $n-1$, hence

$$c = V_n f(A_n) V_n^* b = V_n \tilde{p}(A_n) V_n^* b = \tilde{p}(A)b$$

$$\deg \tilde{p} < n$$

Recall that $f(A) = p(A)$, where p is the (degree $< n$) interpolation polynomial in $\Lambda(A)$.

The eigenvalues of A_n approximate well the outer eigenvalues of A



So Arnoldi constructs a polynomial approximation of f which is accurate on the outer eigenvalues

$$|\tilde{p}(\lambda_i) - f(\lambda_i)| \text{ small for eigenvalues } \lambda_i \text{ "outside"}$$

$$|p(\lambda_k) - f(\lambda_k)| \text{ possibly large for } \lambda_k \in \Lambda(A) \text{ "inside"}$$

(and gives a fast way to evaluate it)

This gives a good approximation of $f(A)$ if the outer eigenvalues are more important than the inner ones.

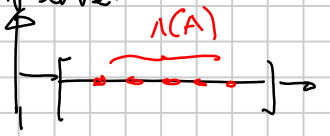
Ex: matrix exponential of diagonalizable $A = V \cdot \Lambda \cdot V^{-1}$

$$\begin{aligned} \exp(A)b &= V \begin{bmatrix} \exp(\lambda_1) \\ \vdots \\ \exp(\lambda_m) \end{bmatrix} \underbrace{V^{-1}b}_{\begin{bmatrix} \sigma_1 \\ \vdots \\ \sigma_m \end{bmatrix}} \\ &= V_1 \exp(\lambda_1) \sigma_1 + V_2 \exp(\lambda_2) \sigma_2 + \dots + V_m \exp(\lambda_m) \sigma_m \end{aligned}$$

The largest summands are usually those with large $\operatorname{Re}(\lambda_i)$, so that $|\exp(\lambda_i)| = \exp(\operatorname{Re}(\lambda_i))$ is large.

Error bound for $c \approx f(A)b$ for Hermitian A

Theorem: Let A be Hermitian, I real interval such that $\lambda(A) \subset [\lambda_{\min}, \lambda_{\max}] \subset I$



Let us consider the best approximation polynomial $s(z)$ to f on I , i.e., the one that achieves minimum

$$\delta = \max_{z \in I} |f(z) - s(z)|$$

over all polynomials of degree $\leq n$.

Then, the approximation provided by Arnoldi $c = V_n f(A_n) V_n^* b$ satisfies

$$\|f(A)b - c\| \leq 2\delta \|b\|.$$

Note that $\|s(A)b - f(A)b\| = \|V \begin{bmatrix} s(\lambda_1) - f(\lambda_1) \\ \vdots \\ s(\lambda_n) - f(\lambda_n) \end{bmatrix} V^* b\|$
 $\leq \|V\| \underbrace{\left\| \begin{bmatrix} s(\lambda_1) - f(\lambda_1) \\ \vdots \\ s(\lambda_n) - f(\lambda_n) \end{bmatrix} \right\|}_{\leq \delta} \underbrace{\|V^*\|}_{=1} \|b\| \leq \delta \cdot \|b\|,$
 Hence the approximation $c = \tilde{f}(A)b$ is within a factor 2 of the best possible degree $\leq n$ polynomial approximation of f .

Proof (of the quasi-optimality theorem):

First, we show that the eigenvalues of A_n are also in I :

$\mu \in A_n$ can be written as a Rayleigh quotient

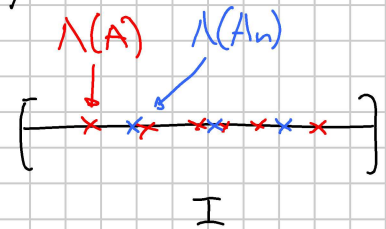
$$\mu = \frac{w^* A_n w}{w^* w} \quad (w \text{ eigenvector of } A_n, A_n w = w \mu)$$

$$= \frac{w^* V_n^* A V_n w}{w^* V_n^* V_n w} \quad (V_n \text{ Arnoldi basis matrix, orthonormal})$$

$$= \frac{y^* A y}{y^* y} = \frac{y^* V \Lambda V^* y}{y^* V V^* y} = \frac{x^* \Lambda x}{x^* x}$$

$$= \frac{\lambda_1 |x_1|^2 + \lambda_2 |x_2|^2 + \dots + \lambda_m |x_m|^2}{|x_1|^2 + \dots + |x_m|^2}$$

= convex combination of the $\lambda: \lambda \in \Lambda(A) \Rightarrow \lambda \in I$.



$$\|f(A)b - c\| = \|f(A)b - V_n f(A) V_n^* b\|$$

$$= \|f(A)b - s(A)b + V_n s(A_n) V_n^* b - V_n f(A_n) V_n^* b\|$$

$$\leq \|f(A)b - s(A)b\| + \|V_n (s(A_n) - f(A_n)) V_n^* b\|$$

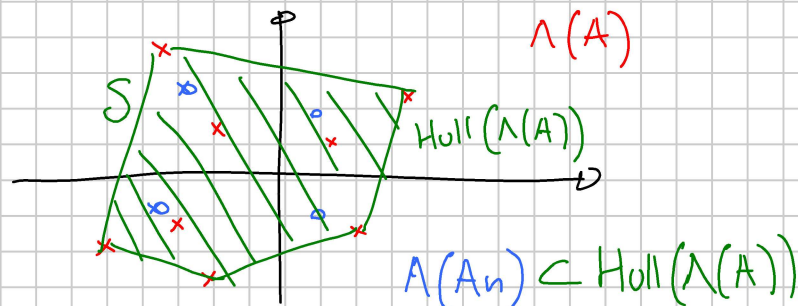
$$\leq \underbrace{\|f(A) - s(A)\|}_{\delta} \cdot \|b\| + \underbrace{\|V_n\|}_{1} \cdot \underbrace{\|s(A_n) - f(A_n)\|}_{\delta} \cdot \underbrace{\|V_n^*\|}_{1} \cdot \|b\|$$

because s satisfies $\delta = \max |f(\lambda) - s(\lambda)|$ because $\lambda(A_n) \in I$, too

$$= 2\delta \|b\|. \quad \square$$

For a normal matrix:

orthogonal V eigenvector matrix, but possibly non-real $\Lambda(A)$



With the same arguments, one can prove the following result:

Let A be normal.

Let $s(z)$ be the best approximation polynomial on $S = \text{Hull}(\lambda(A))$, i.e. the one with minimum

$$\delta = \max_{z \in S} |s(z) - f(z)|$$

over all polynomials of degree $< n$.

$$\text{Then, } \|e - f(A)b\| \leq 2\delta \|b\|$$

What happens if A is not normal?

Def: the numerical range or field of values of $A \in \mathbb{C}^{m \times m}$, possibly non-normal, is

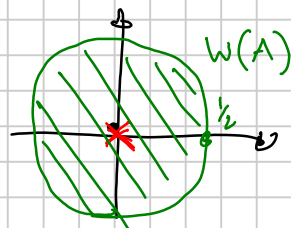
$$W(A) = \left\{ \frac{x^* A x}{x^* x} : \text{for } x \in \mathbb{C}^m, x \neq 0 \right\}$$

(or $x \in$ unit sphere)

$\left\{ \text{sets of Rayleigh quotients of } A \right\}$

For non-normal A , $W(A)$ is bigger than $\text{Hull}(\lambda(A))$.

Ex. $W\left(\begin{bmatrix} 0 & 1 \\ 0 & 0 \end{bmatrix}\right) = \mathcal{B}\left(0, \frac{1}{2}\right)$



Theorem (Grouzeix-Polunin theorem): Let $\gamma = 1 + \sqrt{2} = 2.414\dots$

For any $A \in \mathbb{C}^{m \times m}$ and any f holomorphic on $W(A)$ we have

$$\|f(A)\| \leq \gamma \max_{z \in W(A)} |f(z)|$$

Euclidean induced norm $\|\cdot\|_2$, $\|A\| = \sigma_{\max}(A)$

Open problem (Grouzeix conjecture): This theorem holds also with $\delta=2$.

With this result + the same argument as above:

Theorem: let $A \in \mathbb{C}^{n \times n}$ be any matrix,
 $S(z)$ best approximation polynomial on $W(A)$, then

$$\|c - f(A)b\| \leq 2\delta \gamma \|b\|$$

\uparrow \uparrow constant in the theorem

$$\max_{z \in W(A)} |S(z) - P(z)|$$

Proof: same argument as above.

One can extend the results to so-called rational Arnoldi.

Def:

$$K_{q,n}(A,b) = \left\{ r(A)b, \text{ where } r(z) = \frac{p(z)}{q(z)}, \right. \\ \left. p \text{ any polynomial of degree } < n \right\}$$

There is a variant of Arnoldi that computes a basis of $K_{q,n}(A,b)$, and many of these results continue to hold replacing "polynomial" with "rational function with fixed denominator".