

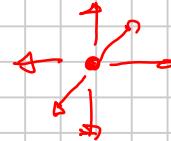
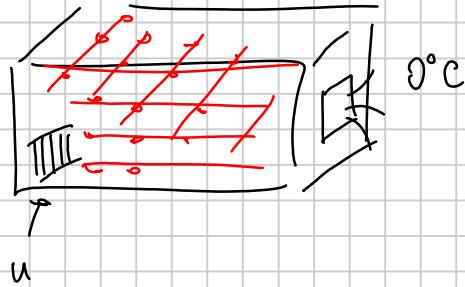
# Large and sparse control systems

Note Title

2023-05-25

Ex.

3D version of heating control problem:



$$A = \begin{pmatrix} -2 & 1 & & \\ 1 & -2 & 1 & \\ & 1 & -2 & 1 \\ & & 1 & -2 \end{pmatrix} \quad \text{in the 1D version,}$$

$$A = \begin{pmatrix} -4 & 1 & 0 & & \\ 1 & -4 & 1 & 0 & \\ 0 & 1 & -4 & 1 & 0 \\ & & 0 & 1 & -4 \\ & & & & \ddots \end{pmatrix} \quad 2D$$

In general,  $A$  large and sparse.

$$\begin{aligned} v &\rightarrow Av \\ v &\rightarrow (A-\alpha I)^{-1}v \end{aligned} \quad \text{the only maps that we will need.}$$

$A \in \mathbb{C}^{n \times n}$  where  $n$  very large

$B \in \mathbb{C}^{n \times m}$  tall thin  $m \ll n$

How do find stabilizability, optimal controls, etc.

In this lecture: just one simple problem,  
solving the Lyapunov equation

$$AX + XA^* + BB^* = 0$$

$$\begin{aligned} A &\in \mathbb{C}^{n \times n}, \quad \Lambda(A) \subset LHP \\ B &\in \mathbb{C}^{n \times 1} \end{aligned}$$

Once we know how to do this,

if  $B$  has dimension  $m \times 1$

$$B = [b_1 \ b_2 \ \dots \ b_m]$$

then we can compute the solution  $X$  to  $AX + XA^* + BB^* = 0$

$$\text{as } X = X_1 + X_2 + \dots + X_m$$

$$X_i \text{ solves } AX_i + X_i A^* + b_i b_i^* = 0$$

since  $BB^* = b_1 b_1^* + b_2 b_2^* + \dots + b_m b_m^*$ , and the equation is linear.

Also, to solve Riccati equations we can use Newton's method once we have an algorithm for Lyapunov eqs.

Problem: even if  $A$  sparse,  $Q = bb^*$  low-rank,

the solution to  $AX + XA^* + bb^* = 0$

is not sparse; how to store it?

(and also if  $(A, b)$  controllable  $X > 0$  has full rank)

Solution: in many practical cases,  $X$  has eigenvalues that decay quickly (exponentially), so we can approximate it with a low-rank matrix.

Lyapunov equation

$$AX + XA^* + bb^* = 0$$

(L)

We wish to convert it to a discrete-time problem and to a Stein equation

$$X - \hat{A} \hat{X} \hat{A}^* = \hat{b} \hat{b}^*$$

(S)

Theorem: Let  $\gamma > 0$ .  $X$  solves (L) if and only if

$X$  solves the Stein equation (S) with

$$\hat{A} = (A - \tau I)^{-1}(A + \tau I), \quad \hat{b} = \sqrt{2\tau} (A - \tau I)^{-1} b$$

Proof: First note that  $\lambda(A - \tau I) \subset \text{LHP}$  if  $\tau > 0$ , so it is invertible

$$(S) \Leftrightarrow X - (A - \tau I)^{-1}(A + \tau I) X (A + \tau I)^* (A - \tau I)^* = 2\tau (A - \tau I)^{-1} b b^* (A - \tau I)^*$$

Multiply by  $(A - \tau I)$  and  $(A - \tau I)^*$  on the left and right

$$\Leftrightarrow (A - \tau I) X (A - \tau I)^* - (A + \tau I) X (A + \tau I)^* = 2\tau b b^*$$

$$\Leftrightarrow -\tau X A^* - A X \tau - \tau X A^* - A X \tau = 2\tau b b^*$$

$$\Leftrightarrow 2\tau (A X + X A^* + b b^*) = 0 \Leftrightarrow (L)$$

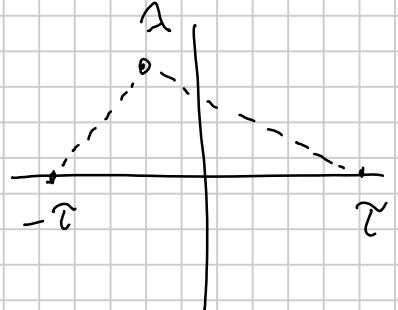
Lemma: if  $\lambda(A) \subset \text{LHP}$ , then  $\lambda(\hat{A}) \subset \text{Unit disc}$

(with  $\hat{A} = (A - \tau I)^{-1}(A + \tau I)$ )

Proof: If  $\{\lambda_1, \lambda_2, \dots, \lambda_n\} = \lambda(A)$ , then

$$\lambda(\hat{A}) = \left\{ \frac{\lambda_i + \tau}{\lambda_i - \tau}, \quad i=1, 2, \dots, n \right\}$$

If  $\lambda \in \text{LHP}$ ,  $\text{dist}(\lambda, -\tau) < \text{dist}(\lambda, \tau)$



hence  $\left| \frac{\lambda + \tau}{\lambda - \tau} \right| < 1$ .

Remark: note that solving  $\hat{x} = \hat{A}x$  with the trapezoidal rule would have returned up to scaling  $x_{k+1} = \hat{A}x_k$

How to solve  $X - \hat{A}X\hat{A}^* = \hat{b}\hat{b}^*$ ?

Idea: fixed point iteration

$$\begin{cases} x_0 = 0 \\ x_{k+1} = \hat{A}x_k \hat{A}^* + \hat{b}\hat{b}^* \end{cases}$$

$$x_1 = \hat{b}\hat{b}^* \quad x_2 = \hat{A}\hat{b}\hat{b}^*\hat{A}^* + \hat{b}\hat{b}^* \quad x_3 = \hat{A}^2\hat{b}\hat{b}^*\hat{A}^* + \hat{A}\hat{b}\hat{b}^*\hat{A}^* + \hat{b}\hat{b}^*$$

$$x_k = \sum_{i=0}^{k-1} \hat{A}^i \hat{b}\hat{b}^* (\hat{A}^*)^i$$

Since  $\Lambda(\hat{A}) \subset \text{unit disc}$ ,  $\hat{A}^i \rightarrow 0$

Lemma: If  $x$  is the exact solution of the Stein equation,

then  $x_k - x = \hat{A}^k(x_0 - x)\hat{A}^*$

Proof: Induction: For  $k=0$  it is trivial, then

$$\begin{aligned} x_{k+1} - x &= \hat{A}x_k \hat{A}^* + \cancel{\hat{b}\hat{b}^*} - (\hat{A}x \hat{A}^* + \cancel{\hat{b}\hat{b}^*}) = \hat{A}(x_k - x)\hat{A}^* \\ &= \hat{A} \stackrel{\text{ind. hypothesis}}{\hat{A}^k} (x_0 - x) \hat{A}^* \hat{A}^* \end{aligned}$$

Hence  $\|x_k - x\| \sim \|\hat{A}\|^{2k} \|x_0 - x\|$ , linear convergence

Low rank formulation:

$$\begin{aligned} x_k &= \hat{b}\hat{b}^* + A\hat{b}\hat{b}^*A^* + A^2\hat{b}\hat{b}^*A^{2*} + \dots + A^{k-1}\hat{b}\hat{b}^*A^{(k-1)*} \\ &= \underbrace{[\hat{b} \quad \hat{A}\hat{b} \quad \hat{A}^2\hat{b} \quad \dots \quad \hat{A}^{k-1}\hat{b}]}_{Z} \cdot \underbrace{[\hat{b} \quad \hat{A}\hat{b} \quad \hat{A}^2\hat{b} \quad \dots \quad \hat{A}^{k-1}\hat{b}]}^* \end{aligned}$$

Algorithm: (ADI iteration with single shift)

Start from  $v = \hat{b}$ ,  $Z = [\hat{b}]$

for  $k=1, 2, 3, \dots$

$$v \leftarrow \hat{A}v = (A - \tau I)^{-1}(A + \tau I)v = (A - \tau I)^{-1}(A - \tau I + 2\tau I)v$$

$$z \leftarrow [z \ v]$$

$$= v + 2\tau(A - I)^{-1}v$$

end

We can return  $\tilde{z}$  s.t.  $X \approx \tilde{z}\tilde{z}^*$

From the convergence speed bound

$$x_k - x = \hat{A}^k(x_0 - x)\hat{A}^{*-k}$$

we get

$$\|x_k - x\| \leq \|\hat{A}^k\|^2 \cdot \|x\|$$

$$\frac{\|x_k - x\|}{\|x\|} \leq \|\hat{A}^k\|^2$$

↓  
 $\rho(\hat{A})^k$

If I stop after  $K$  iterations,

$$\frac{\|x_k - x\|}{\|x\|} \leq \|\hat{A}^k\|^2$$

↓  
 $\rho(\hat{A})^k$

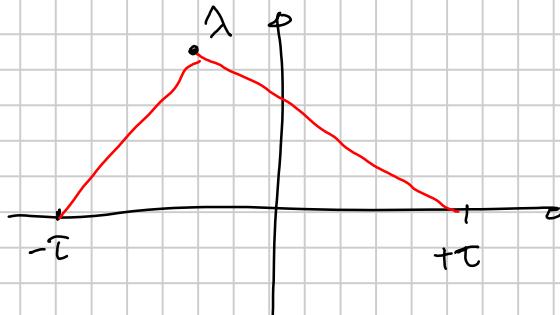
In particular,  $x$  is at distance at most  $\|\hat{A}^k\|^2$   
(relative)

from a matrix  $X_k = z_k z_k^*$  of rank  $K$ .

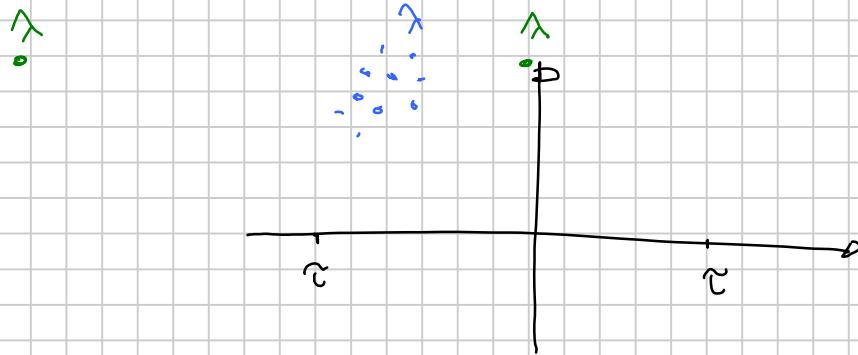
⇒ if  $\rho(\hat{A})$  is small,  $X$  is well-approximated by a low-rank matrix.

What choice of  $\tau$  makes  $\rho(\hat{A})$  as small as possible?

$$\Lambda(\hat{A}) = \left\{ \frac{A + I}{A - \tau} : A \in \Lambda(A) \right\}$$



observation:  $\tau$  should be "around" the eigenvalues of  $A$



If  $\text{Re}(\lambda) \ll \tau$

$$\frac{\lambda - \tau}{\lambda + \tau} \approx 1$$

If  $|\text{Re}(\lambda)| \ll \tau$

$$\frac{\lambda - \tau}{\lambda + \tau} \approx -1$$

In particular, if  $A$  has both very small and large  $|\text{Re}(\lambda)|$  they  $g(\hat{A})$  will be very close to 1, and the algorithm will be very slow

Idea to avoid this: change  $\tau$  at every step!

$$\hat{A}_k = (A - \tau_k I)^{-1} (A + \tau_k I) \quad \hat{b}_k = \sqrt{2\tau} (A - \tau_k I)^{-1} b$$

$$x_{k+1} = \hat{A}_k x_k \hat{A}^* + \hat{b}_k \hat{b}_k^* \quad \text{and different fixed-pt eqn at each step}$$

One can still re-arrange it in a low-rank formulation.

Error bound (proved analogously):

$$x_k - x = \underbrace{\hat{A}_k \hat{A}_{k-1} \cdots \hat{A}_1}_{\in \Lambda(A)} (x_0 - x) \hat{A}_1^* \hat{A}_2^* \hat{A}_3^* \cdots \hat{A}_k^*$$

$$g(\hat{A}_k \hat{A}_{k-1} \cdots \hat{A}_1) = M Q X \frac{|(\lambda + \tau_1)(\lambda + \tau_2) \cdots (\lambda + \tau_k)|}{|\lambda - \tau_1| |\lambda - \tau_2| \cdots |\lambda - \tau_k|}$$

If  $\Lambda(A)$  is composed of only  $K$  distinct real negative

eigenvalues, then after  $K$  steps I can reach

$$g(\hat{A}_k - \hat{A}_0) = 0 \quad \text{by choosing } \tau_i = -\lambda_i \quad (\text{numerical errors zero}).$$

This is an interesting approximation problem:

$$\text{Find } \tau_1, \tau_2, \dots, \tau_K \text{ s.t. } P(x) = (x - \tau_1)(x - \tau_2) \dots (x - \tau_K)$$

is a polynomial of degree  $K$  with

$$|P(\lambda_i)| \text{ small for all } i, \quad |P(-\lambda_i)| \text{ large for all } i$$

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With multiple shifts,

$$\hat{A}_k = (A - \tau_{k1})^{-1}(A + \tau_{k1})$$

$$Z = \begin{bmatrix} \hat{b}_k & \hat{A}_k \hat{b}_{k-1} & \hat{A}_k \hat{A}_{k-1} \hat{b}_{k-2} & \dots & \hat{A}_k \hat{A}_{k-1} \dots \hat{A}_2 \hat{b}_1 \end{bmatrix}$$

In  $Z$  lives in a certain related Krylov subspace:

all the columns are of the form

$$r(\lambda) b, \text{ where } r(x) = \frac{P(x)}{Q(x)} \quad \text{and}$$

$$Q(x) = (x - \tau_k)(x - \tau_{k-1}) \dots (x - \tau_1)$$

This suggests other methods based on Arnoldi: first compute a basis  $V_k$  s.t.  $AV_k = V_k H_{k,k+1}$

then look for  $X \approx V_k Y V_k^*$

$$AX + XA^* + bb^* = 0 \iff$$

$$\iff AV_k Y V_k^* + V_k Y V_k^* A^* + bb^* = 0$$

$$V_k^* AV_k Y + Y V_k^* AV_k + V_k^* b b^* V_k = 0$$

$$A_k Y + Y A_k + \beta e_1 e_1^* \beta^* = 0 \quad \text{a smaller-size Lyapunov equation.}$$